

Exercise Set 6

Exercise 1:

Consider the FRACTIONAL MULTIKNAPSACK problem: Given natural numbers m , n , w_i , c_{ij} , and W_j for $1 \leq i \leq n$ and $1 \leq j \leq m$, find $x_{ij} \in [0, 1]$ satisfying $\sum_{j=1}^m x_{ij} = 1$ for all $1 \leq i \leq n$ and $\sum_{i=1}^n x_{ij} w_{ij} \leq W_j$ for all $1 \leq j \leq m$, such that $\sum_{i=1}^n \sum_{j=1}^m x_{ij} c_{ij}$ is minimum. Provide a polynomial-time combinatorial algorithm for this problem or prove that it is *NP*-hard.

(4 Points)

Exercise 2:

Suppose that in an instance a_1, \dots, a_n of the BIN PACKING problem we have $a_i > \frac{1}{3}$ for $1 \leq i \leq n$.

- (i) Reduce the problem to the CARDINALITY MATCHING problem.
- (ii) Describe a linear-time algorithm that solves the problem.

(3+2 Points)

Exercise 3:

An algorithm for the BIN PACKING problem is called monotone if for inputs S and T with $S \subseteq T$ the algorithm needs at least as many bins for T as for S . Prove:

- (i) Next Fit is monotone.
- (ii) First Fit is not monotone.

(3+3 Points)

Please return the exercises until Tuesday, **May 26th, at 2:15 pm.**