

## Exercise Set 12

### Exercise 1:

Given an instance  $(K_n, c)$  of the TRAVELING SALESMAN PROBLEM and an edge  $e \in E(K_n)$ , show that it is  $NP$ -complete to decide whether  $e$  occurs in an optimal tour.

(5 points)

### Exercise 2:

Why does  $PCP(0, \text{poly}(n)) = NP$  hold with  $\text{poly}(n) := \bigcup_{k \geq 0} n^k$ ?

(3 points)

### Exercise 3:

Consider the following problem: Given  $n$  variables  $x_1, \dots, x_n$  and a set of  $m$  boolean functions  $\Phi = \{\Phi_1, \dots, \Phi_m\}$  using  $k$  variables each, find a truth assignment such that the number of satisfied functions is maximized. The goal is to show that there is no polynomial-time 2-approximation algorithm unless  $P = NP$ .

To do so, find a polynomial reduction  $f$  from SAT to the problem with the following property: If  $I$  is a yes-instance of SAT, then all boolean functions in  $f(I)$  are satisfiable. Otherwise, at most half of the functions in  $f(I)$  are satisfiable. (Hint: Use that there is a  $(\log(n), 1)$ -verifier  $\mathcal{V}$  for SAT.)

(5 points)

Please return the exercises until Tuesday, **July 20th, at 2:15 pm.**