

Have seen: strengthening LP relaxations by adding strong valid inequalities.

E.g.: Knapsack (\rightarrow Carr & Shmoys)
adding flow cut inequalities reduces gap to ≤ 2

in fact: the gap of the strengthened LP is ≈ 2 (see Carr et al. '00)

There is of course a PTAS for max- and min Knapsack. Having an LP with gap ≈ 2 not impressive.

Today: Lift- and Project hierarchies

- strengthening weak LPs systematically

- Idea: (i) find valid relaxation of given problem in higher-dim space "lift"

- (ii) projection to native space is tighter relaxation of problem "project"

There are many Lift & Project hierarchies, we focus on Lasserre that produces SOP relaxations.

Positive Def. Matrices

Symmetric matrix M is positive semidef (psd)

if $x^T M x \geq 0 \quad \forall x$

Equip: (i) any principal submatrix has

$$\det \geq 0$$

(ii) \exists vect. v_i s.t. $M_{ij} = \langle v_i, v_j \rangle$

Round-t Lasserre relax

linear relax of binary problem

Give $K = \{x \in \mathbb{R}^n : Ax \geq b\}$ A $m \times n$ mat.

Introduce new variables $y_I \quad \forall I \subseteq [n], |I| \leq 2t+1$

Intuition think of y_I as representing exact

$$\bigwedge_{i \in I} (x_i = 1)$$

want: $y_\emptyset = 1$ and $y_{\{i\}} = x_i$

$$\boxed{\text{Las}_t(K)} = \left\{ \begin{array}{l} (M_t(y))_{|I|, |J| \leq t} \succeq 0 \\ \left(\sum_{i \in [n]} A_{\ell, i} \cdot y_{I \cup \{i\}} - b_\ell y_I \right)_{|I|, |J| \leq t} \leq 0 \quad \forall \ell \in [m] \\ (M_t^c(y)) \\ y_\emptyset = 1 \end{array} \right.$$

Consistency \downarrow $M_t(y)$ moment matrix

enforces constraints \downarrow $M_t^c(y)$ moment matrix of slacks

$$\text{Las}_t^{\text{proj}}(K) = \{(y_{\{1\}}, \dots, y_{\{n\}}) : y \in \text{Las}_t(K)\}$$

projection of $\text{Las}_t(K)$ onto original space.

Note: set of PSD matrices of certain dim form polyhedral cone

$\Rightarrow \text{Las}_t(K)$ is polyhedron

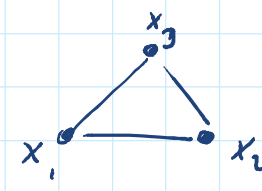
Separation \Rightarrow computing neg. EV
of involved mat if that exists

Can be done in polytime

\Rightarrow can (approx.) optimize over $\text{Las}_t(K)$
in $n^{O(t)}$ mod time.

Ex: Independent set $G = (V, E)$

$$K = \{x \in \mathbb{R}^n : x_u + x_v \leq 1 \quad \forall uv \in E \\ 0 \leq x \leq 1\}$$

$$M_1(\gamma) = \begin{pmatrix} 1 & \gamma_1 & \gamma_2 & \gamma_3 \\ \gamma_1 & 1 & \gamma_{21} & \gamma_3 \\ \gamma_2 & \gamma_{21} & 1 & \gamma_{32} \\ \gamma_3 & \gamma_3 & \gamma_{32} & 1 \end{pmatrix}$$


$$M_1^2 = \begin{pmatrix} 1 - \gamma_1 - \gamma_2 & \gamma_1 - \gamma_1 - \gamma_{12} & \dots \\ \vdots & \gamma_1 - \gamma_1 - \gamma_{12} & \dots \\ \vdots & \vdots & \dots \end{pmatrix}$$

Basic Properties

$$(P1) \quad \text{conv}(K \cap \{0,1\}^n) \subseteq \text{Las}_t^{\text{proj}}(K) \quad \forall t \geq 1$$

give $x \in K \cap \{0,1\}^n$, let

$$\gamma_I = \begin{cases} 1 & : x_i = 1 \quad \forall i \in I \\ 0 & : \text{oth.} \end{cases}$$

note: $(M_\epsilon(y))_{I,J} = y_{I \cup J} = y_I \cdot y_J = (y y^T)_{I,J}$

and $y y^T \succeq 0$

Similarly $M_\epsilon^c(y) \succeq 0 \quad \forall \epsilon$

So: y feas. for $Las_\epsilon(K)$ ✓

(P2) $y \in Las_\epsilon(K)$ Then

(a) $(y_1, \dots, y_n) \in K$

$$M_\epsilon^c(y)_{\emptyset, \emptyset} : \sum_{i=1}^n a_i y_i - b_\emptyset y_\emptyset \stackrel{\text{psd}}{\geq} 0$$

so: \bar{y} sat. all constraints

also: consider principal subm. def.
by \emptyset and I of $M_\epsilon(y)$:

$$\begin{bmatrix} y_\emptyset & y_I \\ y_I & y_I \end{bmatrix} = A$$

$$\det(A) = y_I (1 - y_I) \stackrel{\text{psd}}{\geq} 0$$

$$\Rightarrow 0 \leq y_I \leq 1 \quad \forall I$$

(b) $y_J \leq y_I$ for $I \subseteq J, |I| \leq |J| \leq t$

Consider submatrix for I, J

$$\begin{vmatrix} y_I & y_J \\ y_I & y_J \end{vmatrix} = y_I y_J - y_J^2 = y_J (y_I - y_J) \geq 0$$

$$\Rightarrow y_J \leq y_I$$

Lemma 1 $\eta \in \text{Las}_t(K)$ $i \in [n]$, $0 < b_i < 1$
 $\Rightarrow \eta \in \text{Conv} \{ z \in \text{Las}_{t-1}(K) \mid z_i \in \{0, 1\} \}$

Pf. Define $z_I^{(1)} = \frac{\eta_{I \cup \{i\}}}{b_i}$ $z_I^{(0)} = \frac{\eta_I - b_I \eta_{I \cup \{i\}}}{1 - b_i}$

$$\Rightarrow \eta = b_i z_I^{(1)} + (1 - b_i) z_I^{(0)}$$

$$\text{and } z_i^{(1)} = \frac{b_i}{b_i} = 1, \quad z_i^{(0)} = \frac{b_i - b_i}{1 - b_i} = 0$$

To Do show psd constraints

$$\eta \in \text{Las}_t(K) \Rightarrow M_t(\eta) \succeq 0$$

$$\Rightarrow \exists v_I \text{ s.t. } M_t(\eta)_{I \cup \{i\}} = \langle v_I, v_i \rangle$$

$$\Rightarrow \forall I, J \subseteq [n] \setminus \{i\} : \langle \frac{v_{I \cup \{i\}}}{\sqrt{b_i}}, \frac{v_{J \cup \{i\}}}{\sqrt{b_i}} \rangle = \frac{1}{b_i} \cdot \eta_{I \cup \{i\} \cup J \cup \{i\}} \\ = z_{I \cup J}^{(1)}$$

$$\Rightarrow M_{t-1}(z^{(1)}) \succeq 0$$

$$\text{Also: } \left\langle \frac{v_I - v_{I \cup \{i\}}}{\sqrt{1 - b_i}}, \frac{v_J - v_{J \cup \{i\}}}{\sqrt{1 - b_i}} \right\rangle \\ = \frac{v_I v_J - v_I v_{J \cup \{i\}} - v_{I \cup \{i\}} v_J + v_{I \cup \{i\}} v_{J \cup \{i\}}}{1 - b_i}$$

$$= \frac{\eta_{I \cup J} - b_I \eta_{I \cup J \cup \{i\}}}{1 - b_i} = z_{I \cup J}^{(0)}$$

$$\Rightarrow M_{t-1}(z^{(0)}) \succeq 0$$

Similar for slack matrices. \square

Immediate consequence via induction:

Corollary 1 $y \in \text{Las}_t(K)$, $S \subseteq [n]$, $|S| \leq t$

$$\Rightarrow y \in \text{Conv} \{ z \in \text{Las}_{t-|S|}(K) \mid z_i \in [0,1] \forall i \in S \}$$

Actually can get an explicit convex combination.

Let

$$(Def) \quad y_{\mathcal{I}} = \sum_{K \subseteq \mathcal{I}_0} (-1)^{|\mathcal{I}|} y_{\mathcal{I} \cup \mathcal{I}_0, K}$$

Lemma 2 $y \in \text{Las}_t(K)$, $S \subseteq [n]$

$$\Rightarrow y = \sum_{\substack{\mathcal{I} \cup \mathcal{I}_0 = S \\ y_{\mathcal{I} \cup \mathcal{I}_0, \emptyset} > 0}} y_{\mathcal{I} \cup \mathcal{I}_0, \emptyset} \underbrace{\left(\frac{y_{\mathcal{I} \cup \mathcal{I}_0, \emptyset}}{y_{\mathcal{I} \cup \mathcal{I}_0, \emptyset}} \right)}_{z_{\mathcal{I} \cup \mathcal{I}_0}} \quad (*)$$

EX

check that

$$\sum_{\emptyset} y_{\mathcal{I} \cup \mathcal{I}_0, \emptyset} = 1$$

$$\text{and } z_{\mathcal{I} \cup \mathcal{I}_0} \in \text{Las}_{t-|S|}(K), \text{ and}$$

$$z_i = \begin{cases} 1 & : i \in \mathcal{I}_0 \\ 0 & : i \in \mathcal{I} \end{cases}$$

Proof: fairly straightforward via induction and Lemma 1.

(→ scribble)

Note: Suppose $y \in \text{Las}_n(K)$ then there is a random variable $X \in K \cap \{0,1\}^n$ with

$$(*) \quad \Pr \left[\bigwedge_{i \in \mathcal{I}} X_i = 1 \right] = y_{\mathcal{I}}$$

Simply take decomposition of Lemma 2
and draw $\sum_{i \in \mathcal{I}_0} x_i$ with prob

$$\sum_{i \in \mathcal{I}_0} x_i$$

\Rightarrow $(**)$ holds.

Inclusion-Exclusion:

$$\Pr \left[\bigvee_{i \in \mathcal{I}_0} (X_i = 1) \right] = \sum_{\emptyset \subsetneq K \subseteq \mathcal{I}_0} (-1)^{|K|+1} \Pr \left(\bigwedge_{i \in K} (X_i = 1) \right)$$

$$\begin{aligned} \Rightarrow \Pr \left[\bigwedge_{i \in \mathcal{I}_0} (X_i = 0) \right] &= 1 - \Pr \left[\bigvee_{i \in \mathcal{I}_0} (X_i = 1) \right] \\ &= \sum_{K \subseteq \mathcal{I}_0} (-1)^{|K|} \Pr \left[\bigwedge_{i \in K} (X_i = 1) \right] \end{aligned}$$

Intersecting all events with $\bigwedge_{i \in \mathcal{I}_1} (X_i = 1)$ gives

generalized inclusion-exclusion formula:

$$\begin{aligned} &\Pr \left[\bigwedge_{i \in \mathcal{I}_1} (X_i = 1) \wedge \bigwedge_{i \in \mathcal{I}_0} (X_i = 0) \right] \\ &= \sum_{K \subseteq \mathcal{I}_0} (-1)^{|K|} \Pr \left[\bigwedge_{i \in \mathcal{I}_1 \cup K} X_i = 1 \right] \\ &\stackrel{(**)}{=} \sum_{K \subseteq \mathcal{I}_0} (-1)^{|K|} \sum_{\mathcal{I}_1 \cup K} \\ &\stackrel{(\text{DEF})}{=} \sum_{\emptyset} \end{aligned}$$

Hence:

\sum_{\emptyset}

$$\begin{aligned}
\mathbb{Z}_{\mathcal{I}^0, \mathcal{I}^1} &= \frac{\mathbb{Z}_{\mathcal{I}^0}}{\mathbb{Z}_{\emptyset}} \\
&= \frac{\Pr \left[\bigwedge_{i \in \mathcal{I}^0 \cup \mathcal{I}^1} (X_i = 1) \wedge \bigwedge_{i \in \mathcal{I}^0} (X_i = 0) \right]}{\Pr \left[\bigwedge_{i \in \mathcal{I}^0} (X_i = 0) \wedge \bigwedge_{i \in \mathcal{I}^1} (X_i = 1) \right]} \\
&= \Pr \left[\bigwedge_{i \in \mathcal{I}^1} (X_i = 1) \mid \bigwedge_{i \in \mathcal{I}^0} (X_i = 0) \wedge \bigwedge_{i \in \mathcal{I}^1} (X_i = 1) \right]
\end{aligned}$$

Thus: $\mathbb{Z}_{\mathcal{I}^0, \mathcal{I}^1} \equiv \mathbb{Z}$ conditioned on $\mathcal{I}^0, \mathcal{I}^1$

This global probabilistic view holds even locally:

Lemma 3 $\mathbb{Z} \in \text{Las}_t(k)$ $S \subseteq [n]$, $|S| \leq t$

$\Rightarrow \exists \text{ dist } \mathcal{D}(S) \text{ over } \{0,1\}^{|S|} \text{ with}$

$$\Pr_{z \sim \mathcal{D}(S)} \left[\bigwedge_{i \in I} (z_i = 1) \right] = \mathbb{Z}_I \quad \forall I \subseteq S$$

Pf: again use convex cond. in Lemma 2.

note: nice consistency property:
 $S \neq S' \subseteq [n]$ $|S|, |S'| \leq t$ then

$$\begin{aligned}
\Pr_{z \sim \mathcal{D}(S')} \left(\bigwedge_{i \in I} (z_i = 1) \right) &= \Pr_{z \sim \mathcal{D}(S)} \left(\bigwedge_{i \in I} (z_i = 1) \right) \\
&= \mathbb{Z}_I \quad \forall I \subseteq S \cup S'
\end{aligned}$$

skip?

Ex: 3-colouring

Let $G = (V, E)$ be a 3-colourable graph; i.e.,
 \exists colouring $(\chi_v)_{v \in V}$, $\chi_v \in \{g, r, b\}$

s.t. $\forall uv \in E: x_u \neq x_v.$

$$K = \{x \in [0,1]^{3n} \mid x_{ic} + x_{jc} \leq 1 \quad \forall ij \in E, c \in \{g, r, b\}\}$$

Lemma 4: $\gamma \in \text{Las}_{\geq \epsilon}(K)$

$\Rightarrow \exists$ family of dist. $(\mathcal{D}(S))_{\substack{S \subseteq V \\ |S| \leq \epsilon}}$ s.t.

(a) each event $x \sim \mathcal{D}(S)$ is valid 3-colouring of $G[S]$

(b) $\Pr_{x \sim \mathcal{D}(S)} [x(i_1)=c_1, \dots, x(i_k)=c_k] = \gamma(i_1, c_1, \dots, i_k, c_k)$

$\forall i_1, \dots, i_k \in S, c_1, \dots, c_k \in \{g, r, b\}$

Decomposition Property

Let $K \subseteq \mathbb{R}^n$ be relax of 0,1-polyhedron. We know that $\text{Las}_n(K) = \text{conv}(K \cap [0,1]^n)$. But $\text{Las}_n(K)$ is often stronger.

Ex: Knapsack $K = \{x \in \mathbb{R}_{\geq 0}^n : \sum_{i=1}^n x_i \leq 1.9\}$

v_i : value of object i

Consider $\gamma \in \text{Las}_2(K)$ and let $0 < \gamma_i < 1$.

Claim $\gamma_{\{i,j\}} = 0 \quad \forall i \neq j$

Pf: Suppose not and $\exists i \neq j$ s.t.

$$g_{\{i,j\}} > 0$$

Then pick convex comb. from before:

$$g = g_i g^{(1)} + (1-g_i) g^{(2)}$$

$$\text{know (P2b)}: g_{\{i,j\}}^{(2)} \leq g_i^{(2)} = 0$$

$$\Rightarrow g_{\{i,j\}}^{(1)} > 0 \Rightarrow g_j^{(1)} > 0$$

But this cannot be, since by Lemma 1:
 $\exists g^{(1)}, g^{(2)} \in \text{Las}_0(K) = K$ s.t.

$$g^{(1)} = g_j^{(1)} g^{(11)} + (1-g_j^{(1)}) g^{(12)}$$

and $g^{(11)}, g^{(12)} \in K$ but $g^{(11)}$ viol.
 cap constraint. ▣

Let $\{v_I\}_{\substack{I \subseteq [n] \\ |I| \leq 2}}$ 3-dim vect. s.t.

$$\langle v_I, v_J \rangle = g_{I \cup J} \\ \Rightarrow \langle v_i, v_j \rangle = g_{\{i,j\}} = 0 \quad \forall i \neq j$$

$$\begin{aligned} \Rightarrow \sum_{i=1}^n g_i &= \sum_{i=1}^n \|v_i\|_2^2 = \sum_{i=1}^n \langle v_\emptyset, v_i \rangle = \sum_{i=1}^n \langle v_\emptyset, \frac{v_i}{\|v_i\|_2} \rangle^2 \\ &\leq (\cos \varphi)^2 \|v_\emptyset\|_2^2 \\ &\leq \|v_\emptyset\|_2^2 = 1 \end{aligned}$$

\uparrow
 $\langle v_i, v_i \rangle = \langle v_i, v_\emptyset \rangle$

$K + \mathbb{1}x \leq 1$ defines $\text{conv}(K_n \{2,1\}^n)$.

$$\Rightarrow \text{Las}_2^{P2j}(K) = \text{conv}(K_n \{2,1\}^n)$$

General result:

Lemma 5 $K = \{x \in \mathbb{R}^n : Ax \geq b\}$

Suppose that any $x \in K$ has $\leq t$ ones.

Then

$$\text{Las}_{t+1}^{\text{proj}}(K) = \text{Conv}(K \cap \{0,1\}^n)$$

Pf: Know: $M_{t+1}(y) \geq 0$. Similar argument as before shows that

$$y_I = 0$$

$$\forall I \subseteq [n] : t < |I| \leq 2(t+1) + 1$$

index set dim in $\text{Las}_{t+1}(K)$

Define: $y_I = 0 \quad \forall I, |I| > 2t+3$.

$$M_n(y) = \begin{bmatrix} M_{t+1}(y) & 0 \\ 0 & 0 \end{bmatrix}$$

$\left. \begin{array}{l} |I| \leq t+1 \\ |I| > t+1 \end{array} \right\}$

$\underbrace{\hspace{10em}}_{\substack{|I| \leq t+1 & |I| > t+1}}$

Clearly $M_n(y) \geq 0$ as $M_{t+1}(y) \geq 0$

One can also show that $M_n^c(y) \geq 0 \quad \forall 1 \leq c \leq m$.

$$\Rightarrow \text{Las}_{t+1}(K) \subseteq \text{Las}_n(K) \quad \square$$

Decomposition thm (Karin, Mathieu, Nguyen '11)

Suppose that all $x \in K$ have $\leq k$ ones in $S \subseteq [n]$ and $k \leq t$. Then

$$y \in \text{Las}_S(K) \Rightarrow y \in \text{Conv} \{z : z \in \text{Las}_{S'}(K) : |S'| \leq k\}$$

\cup

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$$z_i \in \{0,1\} \quad \forall i \in S$$